

**1 MARK QUESTIONS**

1. Write the POS form of boolean function  $H$ , which is represented in a truth table as follows:

$X$	$Y$	$Z$	$H$
0	0	0	1
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

2. Write the SOP form of boolean function  $G$ , which is represented in truth table as follows:

$P$	$Q$	$R$	$G$
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	1

3. Write the POS form of boolean function  $H$ , which is represented in a truth table as follows:

$A$	$B$	$C$	$H$
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

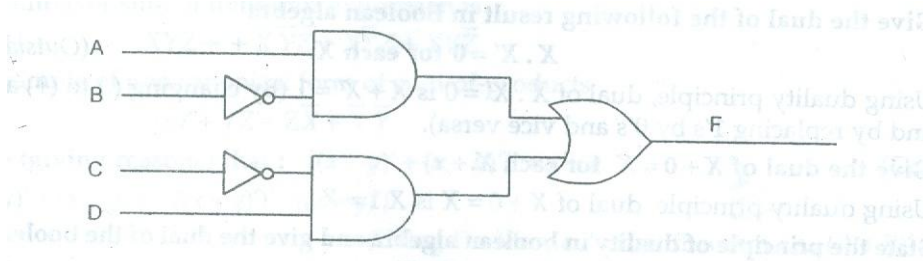
4. Write the POS form a boolean function  $G$ , which is represented in a truth table as follows:

$u$	$v$	$w$	$G$
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	0
1	1	1	1

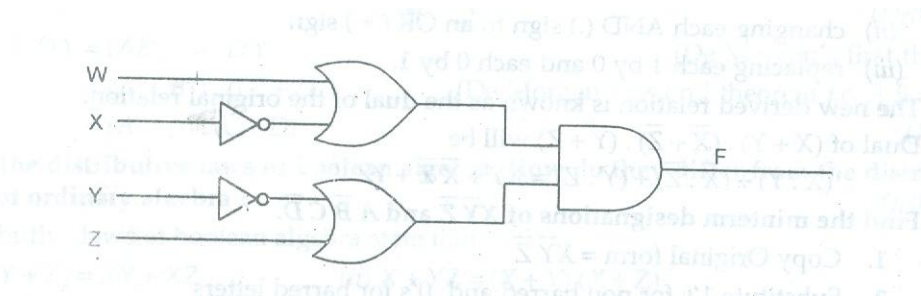
5. Draw a logic circuit diagram for the boolean expression:  $\overline{X}(\overline{Y + Z})$
6. Draw a logic circuit diagram for the boolean expression:  $A.(B+C)$
7. Draw a logic circuit diagram for the boolean expression:  $\overline{A}.(B+C)$
8. Prove that  $X . (X + Y) = X$  by truth table method.
9. Find the complement of the following Boolean function:  

$$F_1 = AB' C'D'$$
10. In the Boolean Algebra, verify using truth table that  $X + XY$  for each  $X, y$  in  $(0, 1)$ .
11. In the Boolean Algebra, verify using truth table that  $(X + Y)' + X'Y'$  for each  $X' Y$  in  $(0, 1)$ .
12. Give the dual of the following result in Boolean Algebra  
 $X . X' =$  for each  $X$ .
13. Define the followings:  
 (a) Minterm      (b) Maxterm    (c) Canonical form

14. Interpret the following logic Circuit as Boolean expression:



15. Interpret the following Logic Circuit as Boolean Expression:

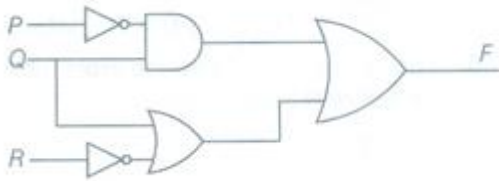


16. Write the dual of the Boolean expression  $A + B' . C$ .
17. Write the dual of the Boolean expression  $(B' + C) . A$ .
18. Represent the boolean expression  $X (Y' + Z)$  with help of NOR gates only.
19. State Demorgan's Laws:
20. Which gates are called Universal gates and why?

## 2 Marks Questions

- 1 State DeMorgan's laws. Verify one of the DeMorgan's laws using a truth table.
- 2 Draw a logic circuit for the following boolean expression.  

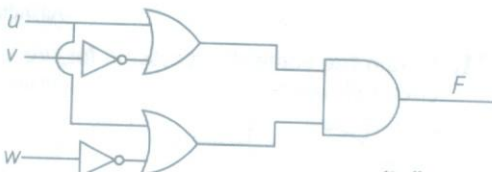
$$A.B' +(C+B') .A'$$
- 3 Obtain the Boolean expression for the logic circuit shown below:



4 Verify the following using boolean expression using truth table:

- (i)  $X + 0 = X$
- (ii)  $X + X' = 1$

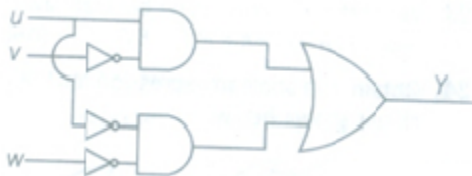
5 Write the equivalent Boolean expression for the following logic circuit:



6 Verify the following boolean expression using truth table:

- (i)  $X \cdot X' = 0$
- (ii)  $X + 1 = 1$

7 Write the equivalent boolean expression for the following logic circuit:



- 8. Represent the Boolean expression  $X \cdot Y' + Z$  with the help of NOR gates only.
- 9. Represent the Boolean expression  $(X + Y) \cdot Z$  with the help of NAND gates only.

### 3 Marks Questions

1. Obtain the minimal form for the following boolean expression using Karnaugh's Map:  
 $F(A, B, C, D) = \sum(1, 4, 5, 9, 11, 12, 13, 15)$
2. Obtain a simplified form for the following boolean expression using Karnaugh's Map:  
 $F(P, Q, R, S) = \sum(0, 2, 4, 5, 6, 7, 8, 10, 13, 15)$
3. Obtain the minimal form for the following boolean expression using Karnaugh's Map:  
 $H(P, Q, R, S) = \sum(0, 1, 2, 3, 5, 7, 8, 9, 10, 14, 15)$
4. Obtain the minimal form for the following boolean expression using Karnaugh's Map:  
 $F(U, V, W, Z) = \sum(0, 1, 2, 3, 6, 7, 8, 9, 10, 13, 15)$
5. Reduce the following boolean expression using K-map:  
 $F(P, Q, R, S) = \sum(1, 2, 3, 4, 5, 6, 7, 8, 10)$
6. Reduce the following boolean expression using K-map:  
 $F(A, B, C, D) = \sum(2, 3, 4, 5, 6, 7, 8, 10, 11)$

7. Reduce the following boolean expression using K-map:  
 $F(A, B, C, D) = \sum(0, 1, 2, 4, 5, 6, 8, 10,)$
8. Reduce the following boolean expression using K-map:  
 $F(P, Q, R, S) = \sum(0, 1, 2, 4, 5, 6, 8, 12)$
9. Reduce the following boolean expression using K-map:  
 $F(A, B, C, D) = \sum(3, 4, 5, 6, 7, 13, 15)$
10. Reduce the following boolean expression using K-map:  
 $F(u, v, w, z) = \sum(3, 5, 7, 10, 11, 13, 15)$

## BOOLEAN ALGEBRA-SOLUTION

### 1 MARK QUESTIONS

1.	<i>X</i>	<i>Y</i>	<i>Z</i>	<i>H</i>	Maxterm
	0	0	0	1	$X + Y + \underline{Z}$
	0	0	1	0	$X + \underline{Y} + Z$
	0	1	0	1	$X + \underline{Y} + \underline{Z}$
	0	1	1	1	$\underline{X} + Y + Z$
	1	0	0	1	$\underline{X} + Y + \underline{Z}$
	1	0	1	0	$\underline{X} + \underline{Y} + Z$
	1	1	0	0	$\underline{X} + \underline{Y} + \underline{Z}$
	1	1	1	1	$X + \underline{Y} + \underline{Z}$

To get the POS form, we need to maxterms for all those input combinations that produce output as

0. Thus,

$$H(X, Y, Z) = (X + Y + \underline{Z}) \cdot (\underline{X} + Y + \underline{Z}) \cdot (\underline{X} + \underline{Y} + \underline{Z})$$

2.	<i>P</i>	<i>Q</i>	<i>R</i>	GMinterm
	0	0	0	$\underline{P} \cdot \underline{Q} \cdot \underline{R}$
	0	0	1	$\underline{P} \cdot \underline{Q} \cdot R$
	0	1	0	$\underline{P} \cdot Q \cdot \underline{R}$
	0	1	1	$\underline{P} \cdot Q \cdot R$
	1	0	0	$P \cdot \underline{Q} \cdot \underline{R}$
	1	0	1	$P \cdot \underline{Q} \cdot R$
	1	1	0	$P \cdot Q \cdot \underline{R}$
	1	1	1	$P \cdot Q \cdot R$

To get the SOP form, we need to sum minterms for all those combinations that produce output as 1. Thus,

$$G(P, Q, R) = (\underline{P} \cdot \underline{Q} \cdot \underline{R}) + (\underline{P} \cdot \underline{Q} \cdot R) + (P \cdot \underline{Q} \cdot \underline{R}) + (\underline{P} \cdot Q \cdot R) + (P \cdot Q \cdot R)$$

3.

	<i>A</i>	<i>B</i>	<i>C</i>	<i>H</i>	Maxterm
	0	0	0	0	$A + B + C$
	0	0	1	1	$A + B + \underline{C}$
	0	1	0	1	$A + \underline{B} + C$
	0	1	1	1	$A + B + C$
	1	0	0	1	$\underline{A} + B + C$
	1	0	1	0	$\underline{A} + B + \underline{C}$
	1	1	0	0	$\underline{A} + \underline{B} + C$
	1	1	1	1	$\underline{A} + \underline{B} + \underline{C}$

To get the POS from, we need to product maxterms for all those input combinations that produce output as 0. Thus,

$$H(A, B, C) = (A + B + C) \cdot (\overline{A} + B + \overline{C}) \cdot (\overline{A} + \overline{B} + C).$$

4.

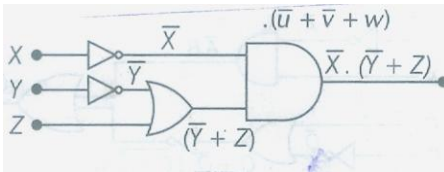
$u$	$v$	$w$	<b>GMaxterm</b>
0	0	0	1 $u + v + w$
0	0	1	1 $u + v + \overline{w}$
0	1	0	0 $u + \overline{v} + w$
0	1	1	0 $u + v + w$
1	0	0	1 $\overline{u} + v + w$
1	0	1	1 $\overline{u} + \overline{v} + w$
1	1	0	0 $\overline{u} + v + \overline{w}$
1	1	1	1 $\overline{u} + \overline{v} + \overline{w}$

as

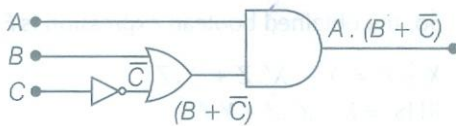
To get the POS form, we need to product maxterms for all those input combinations that produce output 0. Thus,

$$G(u, v, w) = (u + v + w) \cdot (u + v + \overline{w}) \cdot (\overline{u} + v + w) \cdot (\overline{u} + \overline{v} + w)$$

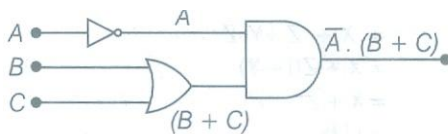
5.



6.



7.



8.

X	Y	X + Y	X . (X + Y)
0	0	0	0
0	1	1	0
1	0	1	1
1	1	1	1

From the above table it is obvious that  $X \cdot (X + Y) = X$  because both the columns are identical.

9.

$$(AB' + CD') = (AB')' \cdot (C'D')'$$

$$= (A' + B'') \cdot (C'' + D'')$$

(De Morgan's first theorem) —  
 (De Morgan's second theorem i.e.  $A \cdot B = \overline{\overline{A} + \overline{B}}$ )

B)

$$= (A' + B) . (C + D) \quad (X'' = X)$$

10. As the expression  $X + XY$  is a two variable expression, so we require possible combinations of values of X, Y. Truth Table will be as follows:

X	Y	$X + Y$	$X . (X + Y)$
0	0	0	0
0	1	0	0
1	0	0	1
1	1	1	1

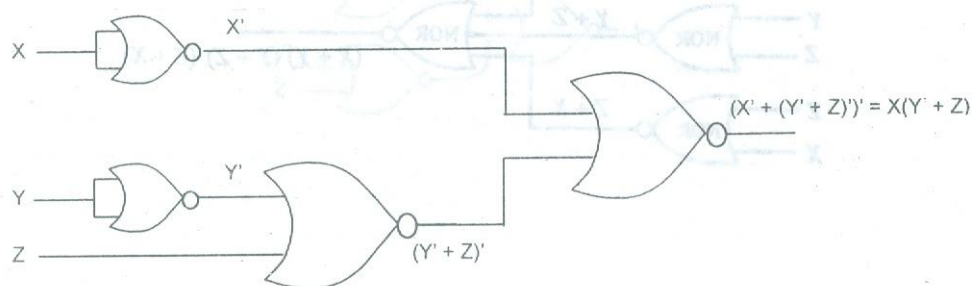
Comparing the columns  $X + XY$  and  $X$ , we find, contents of both the columns are identical, hence verified.

11. As it is a 2 variable expression, truth table will be as follows:

X	Y	$X + Y$	$(X + Y)'$	$X'$	$Y'$	$X'Y'$
0	0	0	1	1	1	1
0	1	1	0	1	0	0
1	0	1	0	0	1	0
1	1	1	0	0	0	0

Comparing the columns  $(X + Y)'$  and  $X'Y'$ , both of the columns are identical, hence verified.

12. Using duality principle, dual of  $X . X' = 0$  is  $X + X' = 1$  (By changing  $(.)$  to  $(+)$  and viceversa and by replacing 1's by 0's and vice versa).
13. (a) A Minterm is a product of all the literals (with or without the bar) within the logic system.  
 (b) A Maxterm is a sum of all the literals (with or without the bar) within the logic system.  
 (c) A Boolean expression composed entirely either of minterms or Maxterms is referred to as canonical expression.
14.  $F = \overline{AB} + \overline{CD}$
15.  $F = \overline{(W + X)} (Y + Z)$ .
16. Dual of the Boolean expression  $A + B' . C$  is  $A . (B' + C)$ .
17. Dual of the Boolean expression  $(B' + C) . A$  is  $(B' . C) + A$ .
18. The given expression may also be written as



19. De Morgan's first theorem. It states that  $X + Y = \overline{X \cdot Y}$   
 De Morgan's second theorem. It states that  $X \cdot Y = \overline{\overline{X} + \overline{Y}}$

20. NAND and NOR gates are less expensive and easier to design. Also, other switching functions and (AND, OR) can easily be implemented using NAND/NOR gates. Thus, these (NAND/NOR) gates are also referred to as *Universal Gates*.

## 2 Marks Questions

1. DeMorgan's Laws:

It states that

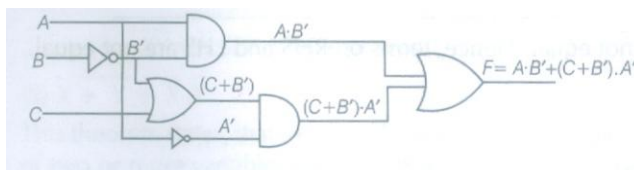
(i)  $\overline{A + B} = \overline{A} \cdot \overline{B}$       (ii)  $\overline{A \cdot B} = \overline{A} + \overline{B}$

Truth table for  $\overline{A + B} = \overline{A} \cdot \overline{B}$

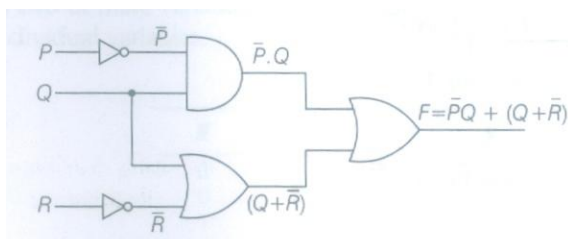
A	B	A+B	$\overline{A+B}$	$\overline{A}$	$\overline{B}$	$\overline{A \cdot B}$
0	0	0	1	1	1	1
0	1	1	0	1	0	0
1	0	1	0	0	1	0
1	1	1	0	0	0	0

Column 4 and Column 7 are equal, first law is proved.

2.  $F = A \cdot B' + (C + B') \cdot A'$



3.





So, the obtained boolean expression is  $F = P'Q + (Q + R')$

4. (i)  $X + 0 = X$

X	0	X + 0
0	0	0
1	0	1

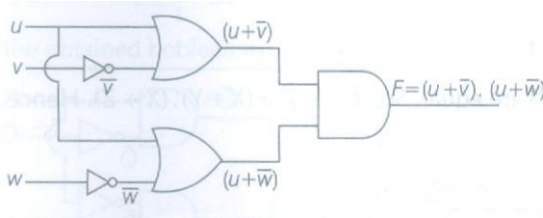
So,  $X + 0 = X$

(ii)  $X + X' = 1$

X	X'	X + X'
0	1	1
1	0	1

As  $X + X' = 1$ . Hence proved.

5.



So, the obtained boolean expression if  $F = (u + \bar{v}).(u + \bar{w})$

6.

(i)  $X.X' = 0$

X	X'	X.X'
0	1	0
1	0	0

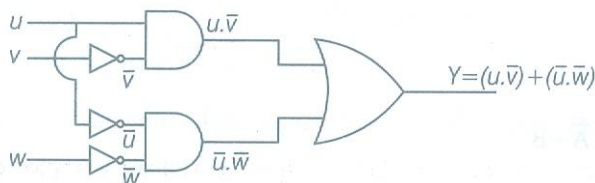
As  $X.X' = 0$ . Hence proved.

(ii)  $X + 1 = 1$

X	1	X + 1
0	1	1
1	1	1

As  $X + 1 = 1$ . Hence proved.

7.



So, the obtained boolean expression is  $Y = (u.\bar{v}) + (\bar{u}.\bar{w})$

law)

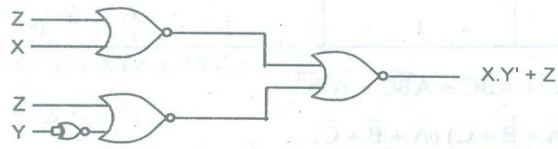
8.  $X.Y' + Z$

$$= Z + XY'$$

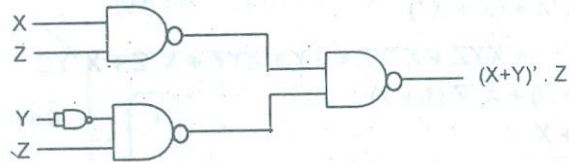
$$= (Z + X)(Z + Y')$$

$$[X + Y = Y + X]$$

$$[X + YZ = (X + Y)(X + Z)]$$



9.  $(X + Y)'. Z = X'. Z + Y'. Z$



### 3 Marks Questions

1.  $F(A, B, C, D) = \Sigma(1, 4, 5, 9, 11, 12, 13, 15)$

AB \ CD	C'D'	C'D	CD	CD'
A'B'	0	1	3	2
A'B	4	5	7	6
AB	12	13	15	14
AB'	8	9	11	10

There are 3 Quads:

Quad 1 ( $m_1 + m_5 + m_9 + m_{13}$ ) reduces to  $C'D$

Quad 2 ( $m_4 + m_5 + m_{12} + m_{13}$ ) reduces to  $BC'$

Quad 3 ( $m_9 + m_{11} + m_{13} + m_{15}$ ) reduces to  $AD$

Hence, the final expression is:

$$F(A, B, C) = C'D + BC' + AD$$

2.  $F(P, Q, R, S) = \Sigma(0, 2, 4, 5, 6, 7, 8, 10, 13, 15)$

PQ \ RS	$\bar{R}\bar{S}$	$\bar{R}S$	$RS$	$R\bar{S}$
$\bar{P}\bar{Q}$	0	1	3	2
$\bar{P}Q$	4	5	7	6
$PQ$	12	13	15	14
$P\bar{Q}$	8	9	11	10

There are three Quads:

Quad 1 ( $m_4 + m_5 + m_6 + m_7$ ) reduces to  $\bar{P}Q$

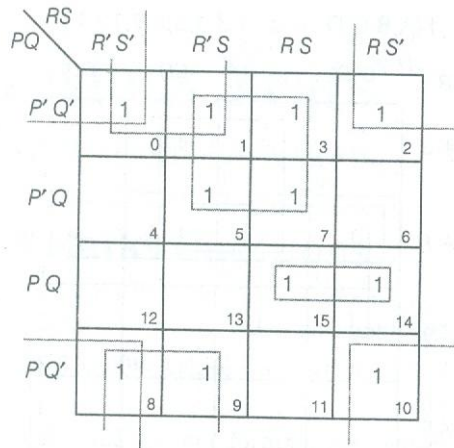
Quad 2 ( $m_5 + m_7 + m_{13} + m_{15}$ ) reduces to  $QS$

Quad 3 ( $m_0 + m_2 + m_8 + m_{10}$ ) reduces to  $QS$

Hence, the final expression is:

$$F(P, Q, R, S) = \bar{P}Q + QS + QS$$

3.  $H(P, Q, R, S) = \Sigma(0, 1, 2, 3, 5, 7, 8, 9, 10, 14, 15)$



There are three Quads and 1 Pair:

Quad 1 ( $m_0 + m_2 + m_8 + m_{10}$ ) reduces to  $\overline{Q}S$

Quad 2 ( $m_1 + m_3 + m_5 + m_7$ ) reduces to  $PS$

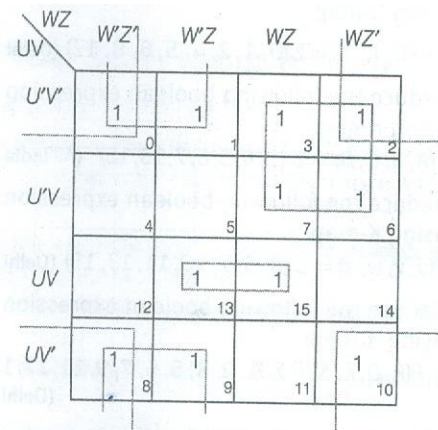
Quad 3 ( $m_0 + m_1 + m_8 + m_9$ ) reduces to  $QR$

Pair 1 ( $m_{14} + m_{15}$ ) reduces to  $PQR$

Hence, the final expression is:

$$H(P, Q, R, S) = \overline{Q} S + P S + Q R + PQR$$

4.  $F(U, V, W, Z) = \Sigma(0, 1, 2, 3, 6, 7, 8, 9, 10, 13, 15)$



There are three Quads and 1 Pair:

Quad 1 ( $m_0 + m_2 + m_8 + m_{10}$ ) reduces to  $VZ$

Quad 2 ( $m_0 + m_1 + m_8 + m_9$ ) reduces to  $VW$

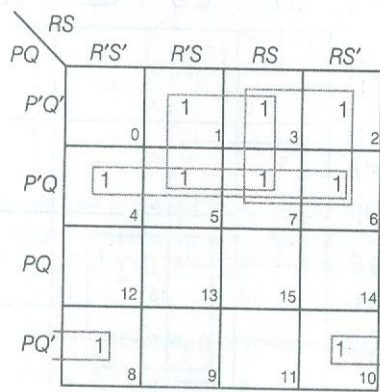
Quad 3 ( $m_2 + m_3 + m_6 + m_7$ ) reduces to  $UW$

Pair 1 ( $m_{13} + m_{15}$ ) reduces to  $UVZ$

Hence, the final expression is:

$$F(U, V, W, Z) = VZ + VW + UW + UVZ$$

5.  $F(P, Q, R, S) = \Sigma(1, 2, 3, 4, 5, 6, 7, 8, 10)$



There are three Quads and 1 Pair:

Quad 1 ( $m_1 + m_3 + m_5 + m_7$ ) reduces to  $P'S$

Quad 2 ( $m_4 + m_5 + m_6 + m_7$ ) reduces to  $P'Q$

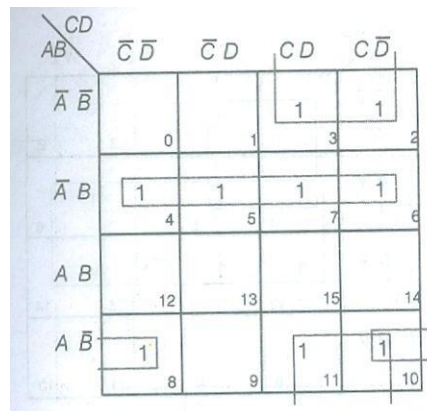
Quad 3 ( $m_2 + m_3 + m_6 + m_7$ ) reduces to  $P'R$

Pair 1 ( $m_8 + m_{10}$ ) reduces to  $PQ'S'$

Hence, the final expression is:

$$F(P,Q,R,S) = P'S + P'Q + P'R + PQ'S'$$

6.  $F(A, B, C, D) = \Sigma(1, 3, 4, 5, 6, 7, 12, 13)$



There are 2 Quads and 1 Pair:

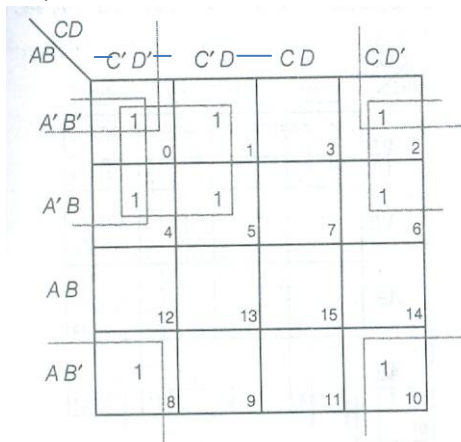
Quad 1 ( $m_4 + m_5 + m_6 + m_7$ ) reduces to  $\overline{A}B$

Quad 2 ( $m_2 + m_3 + m_{10} + m_{11}$ ) reduces to  $B\overline{C}$

Pair 1 ( $m_8 + m_{10}$ ) reduces to  $ABD$

$$F(A, B, C, D) = AB$$

7.  $F(A, B, C, D) =$



Hence, the final expression is:

$$+BC + ABD$$

$$\Sigma(1, 2, 4, 5, 6, 8, 10)$$

There are 3 Quads:

Quad 1 ( $m_0 + m_1 + m_4 + m_5$ ) reduces to  $A'C'$

Quad 2 ( $m_0 + m_2 + m_8 + m_{10}$ ) reduces to  $B'D'$

Quad 3 ( $m_0 + m_2 + m_4 + m_6$ ) reduces to  $A'D'$

Hence, the final expression is:

$$F(A, B, C, D) = A'C' + B'D' + A'D'$$

8.  $F(P, Q, R, S) = \Sigma(1, 2, 4, 5, 6, 8, 10, 12)$

RS PQ	$\bar{R}\bar{S}$	$\bar{R}S$	$RS$	$R\bar{S}$
$\bar{P}\bar{Q}$	1	1		1
$\bar{P}Q$	1	1		1
$PQ$	1			
$P\bar{Q}$	1			

There are 3 Quads:

Quad 1 ( $m_0 + m_4 + m_8 + m_{12}$ ) reduces  $\bar{R}\bar{S}$

Quad 2 ( $m_0 + m_1 + m_4 + m_5$ ) reduces  $PS$

Quad 3 ( $m_0 + m_2 + m_4 + m_6$ ) reduces  $RS$

Hence, the final expression is:

$$F(P, Q, R, S) = RS + PS + \bar{R}\bar{S}$$

9.  $F(A, B, C, D) = \Sigma(3, 4, 5, 6, 7, 13, 15)$

CD AB	$\bar{C}\bar{D}$	$\bar{C}D$	$CD$	$C\bar{D}$
$\bar{A}\bar{B}$			1	
$\bar{A}B$	1	1	1	1
$AB$		1	1	
$A\bar{B}$				

There are 2 Quads and 1 Pair:

Quad 1 ( $m_4 + m_5 + m_6 + m_7$ ) reduces to  $AB$

Quad 2 ( $m_5 + m_7 + m_{13} + m_{15}$ ) reduces to  $BD$

Pair 1 ( $m_3 + m_7$ ) reduces to  $ACD$

Hence, the final expression is:

$$F(A, B, C, D) = AB + BD + ACD$$

10.  $F(u, v, w, z) = \Sigma (3, 5, 7, 10, 11, 13, 15)$

		wz			
		$\bar{w}\bar{z}$	$\bar{w}z$	$wz$	$w\bar{z}$
uv	$\bar{u}\bar{v}$	0	1	3	2
	$\bar{u}v$	4	5	7	6
uv	$\bar{u}v$	12	13	15	14
	$u\bar{v}$	8	9	11	10

There are 2 Quads and 1 Pair:

Quad 1 ( $m_3 + m_7 + m_{11} + m_{15}$ ) reduces to  $wz$

Quad 2 ( $m_5 + m_7 + m_{13} + m_{15}$ ) reduces to  $vz$

Pair 1 ( $m_{10} + m_{11}$ ) reduces to  $uvw$

Hence, the final expression is:

$$F(u, v, w, z) = wz + vz + uvw$$